

Efficient Local Resorting Techniques with Space Filling Curves

Applied to a Parallel Tsunami Simulation Model

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Alfred Wegener Institute for Polar and Marine Research
Bremerhaven, 22 - 25 August 2011

Outline

- introducing TsunAWI
- motivation for resorting
- construction of Hilbert space filling curve (SFC) ordering
- comparison to other sortings
- conclusions

The AWI Tsunami Modell TsunAWI

TsunAWI in a nutshell

- shallow water equations with inundation
- unstructured $P_1 - P_1^{\text{NC}}$ finite element grid
- explicit time stepping scheme
- OpenMP parallel Fortran90 code

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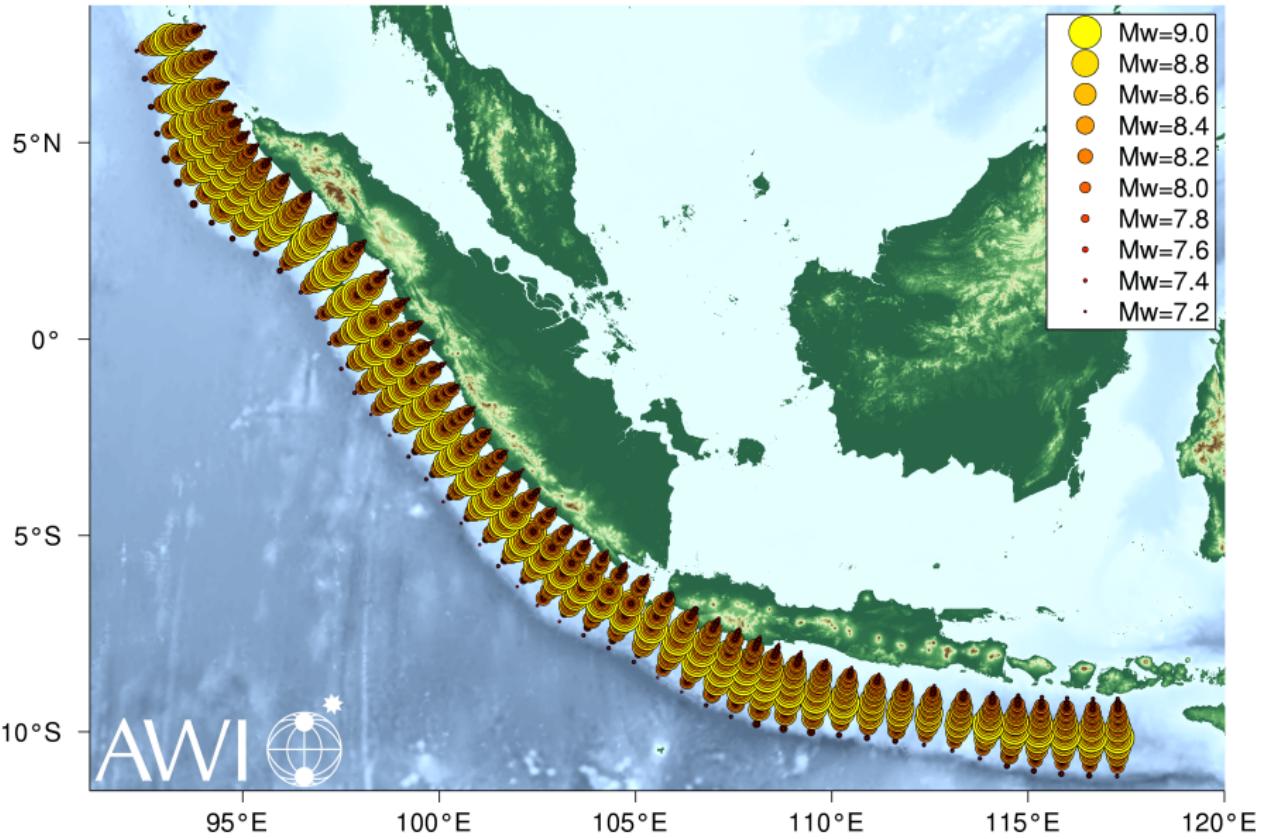
Most important application:

German-Indonesian Tsunami Early Warning System

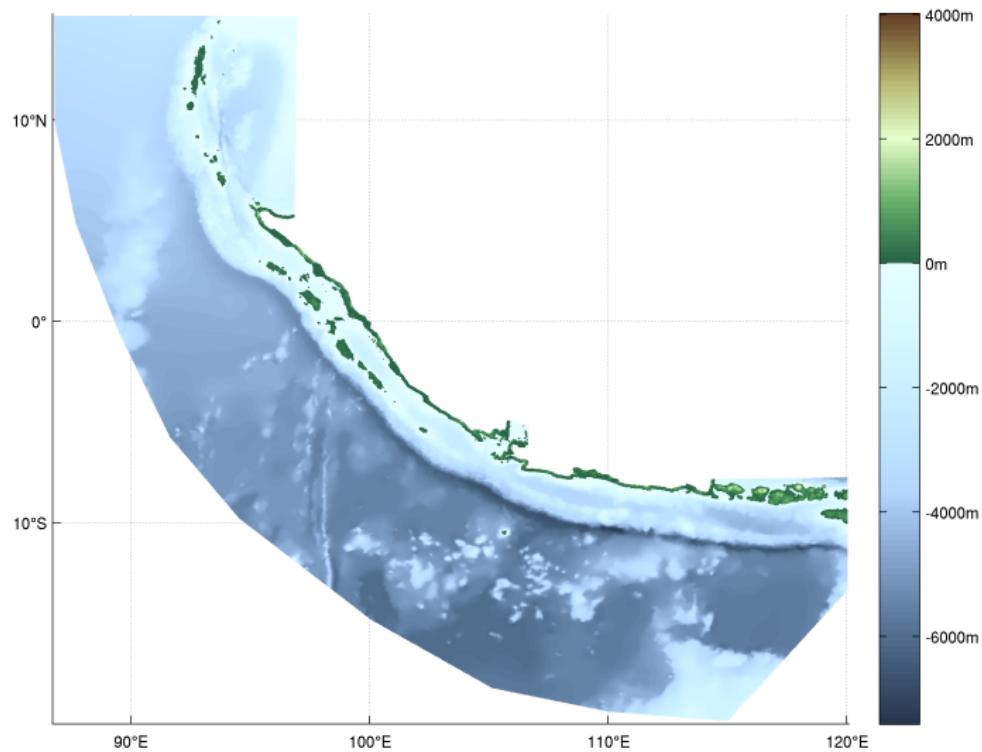


- 3470 scenarios for different prototypic ruptures
- 3h modeltime (10.800 timesteps of 1s)

TsunAWI Scenario Repository for GITEWS, March 2011
scenarios for 3470 prototypic ruptures (RuptGen2.1)

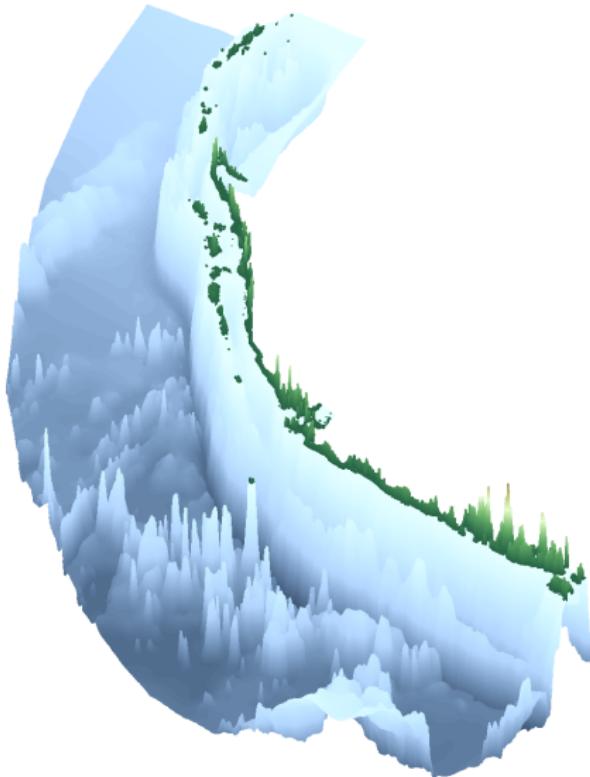


TsunAWI: example for a computational domain regional grid for the Sunda Arc



TsunAWI: example for a computational domain

regional grid for the Sunda Arc

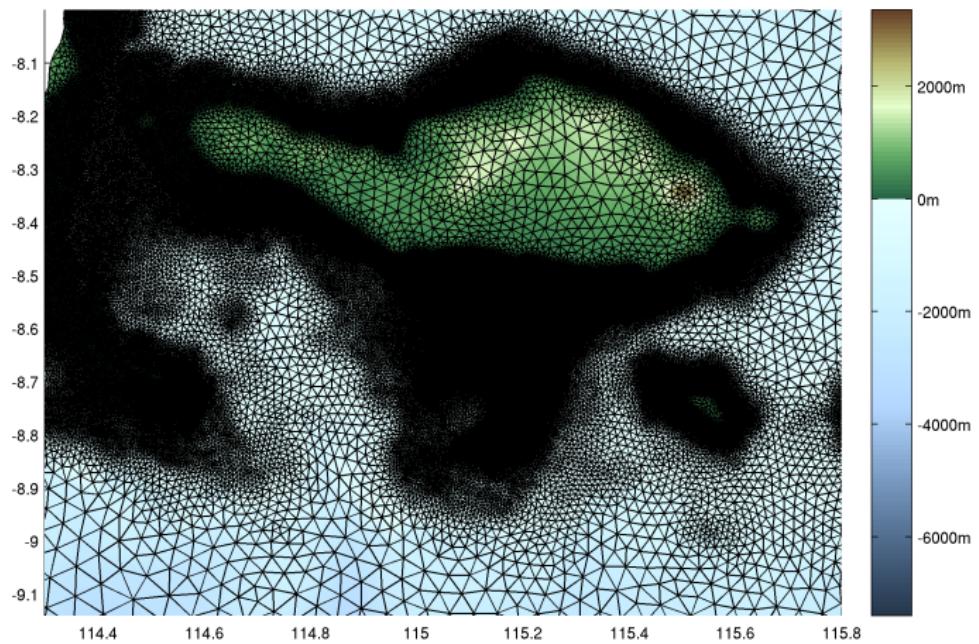


The computational grid discretizes the domain with

- varying resolution
 - 50m areas of interest
 - 500m all other coastal areas
 - 15km deep ocean
- 2.366.319 nodes
- 4.721.884 elements

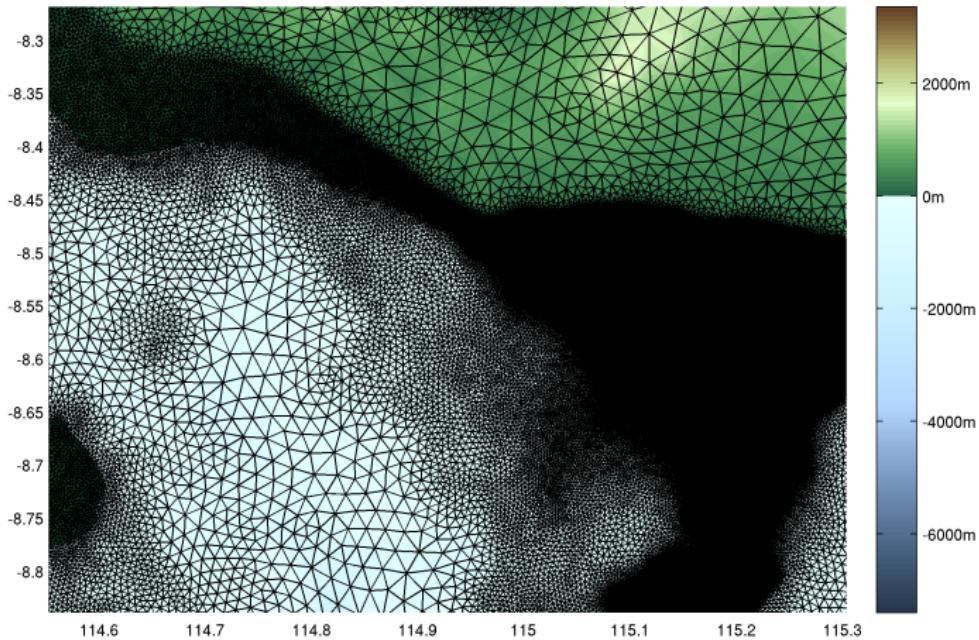
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regional grid for the Sunda Arc, focus on Bali



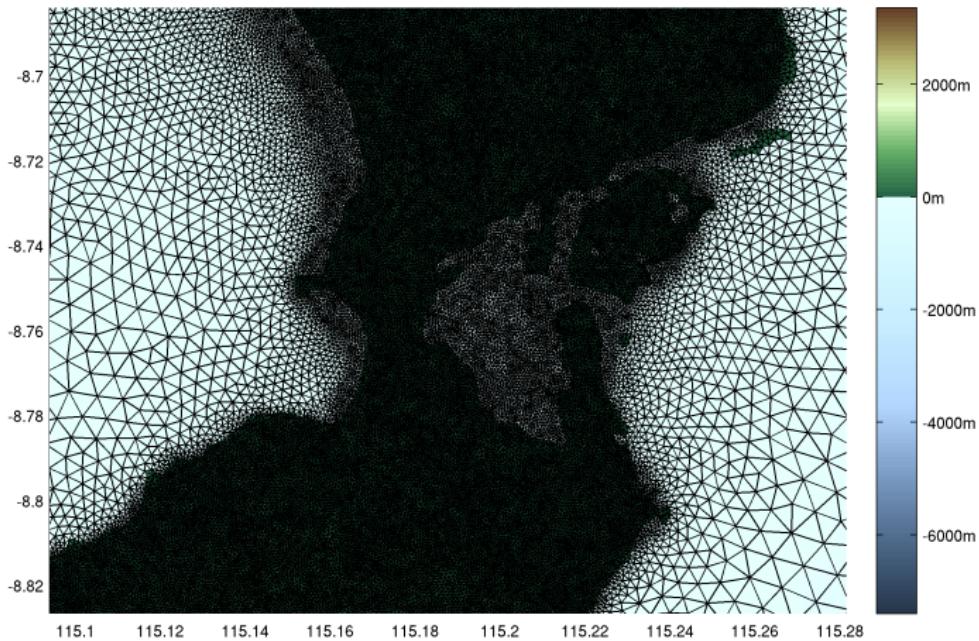
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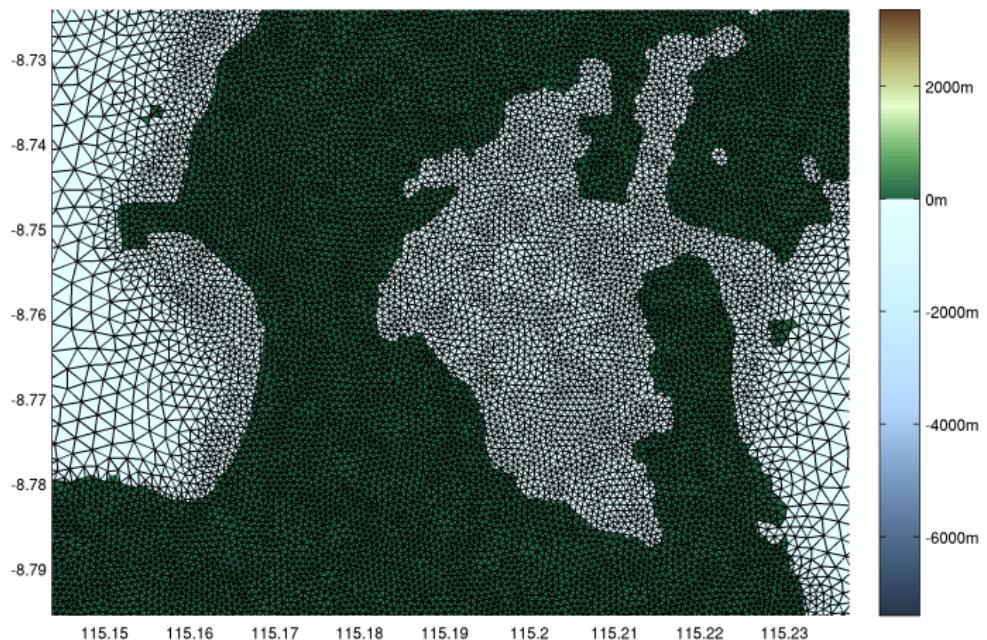
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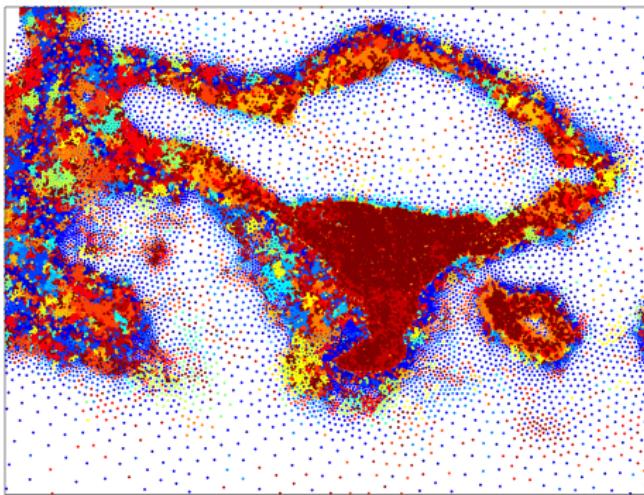
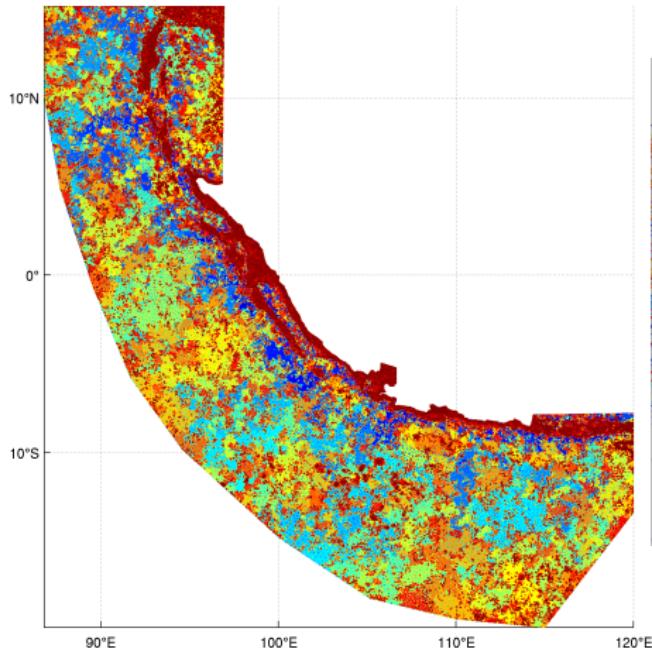
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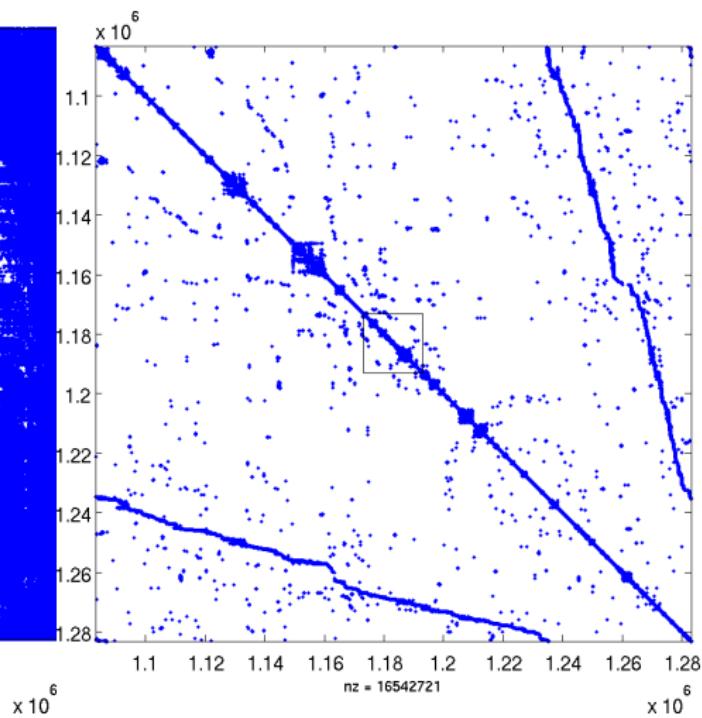
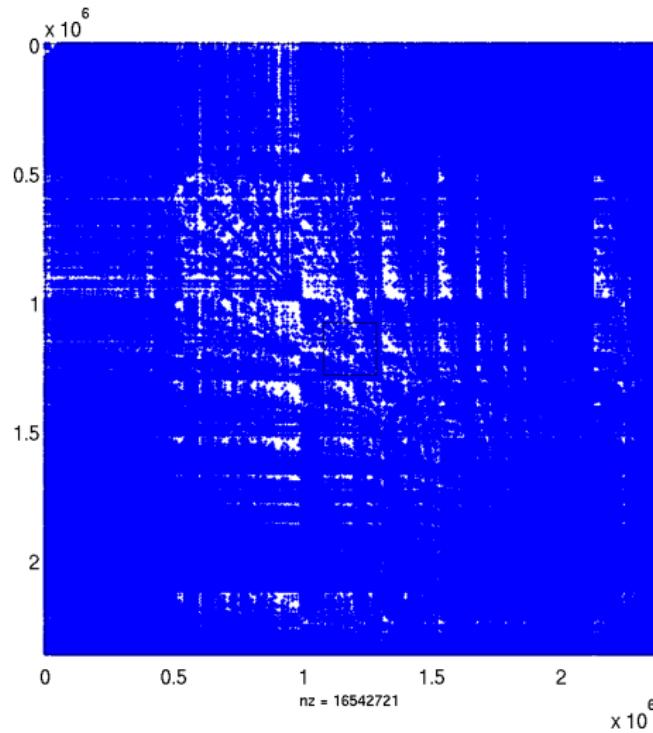


TsunAWI: example for a computational domain

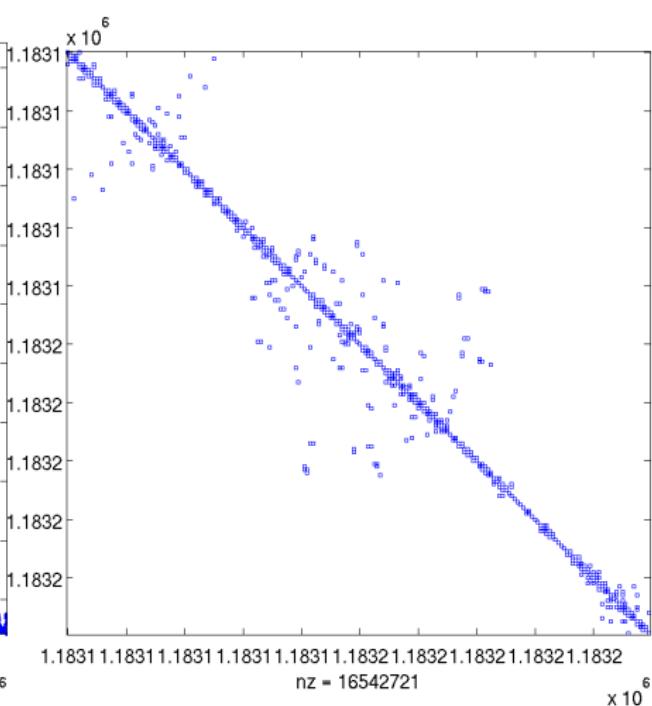
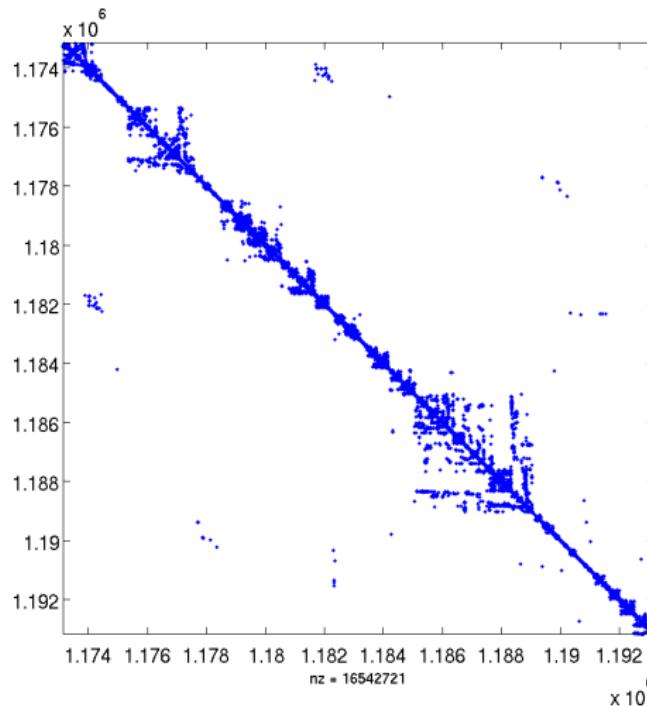
Original numbering of nodes as provided by the grid generator



adjacency matrix, original grid



adjacency matrix, original grid



Motivation for resorting

Data locality on the original grid is **very, very** bad.

E.g., each computation on all nodes of one element results in at least one cache miss.

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Most time consuming routines in every timestep:

`compute_velocity_at_nodes` $v(\text{node}) = F(\text{adjacent edges, elems})$

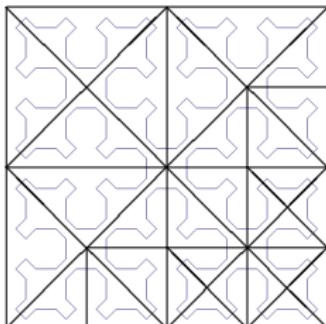
`compute_velocity` $v(\text{edge}) = F(\text{adjacent elems, nodes})$

`compute_ssh` $\text{ssh}(\text{node}) = F(\text{adjacent elems, nodes})$

`compute_gradient` $\text{grad}_{x,y}(\text{elem}) = F(\text{adjacent nodes})$

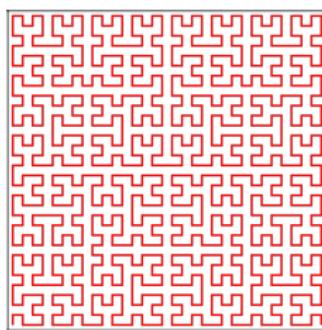
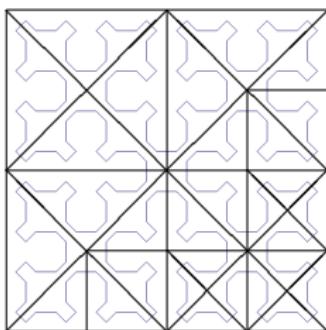
Ideas for resorting

- SFC like Sierpinski curve in adaptive grid (J. Behrens et al., KlimaCampus Uni Hamburg) could help.
But how to derive SFC for highly unstructured grid?



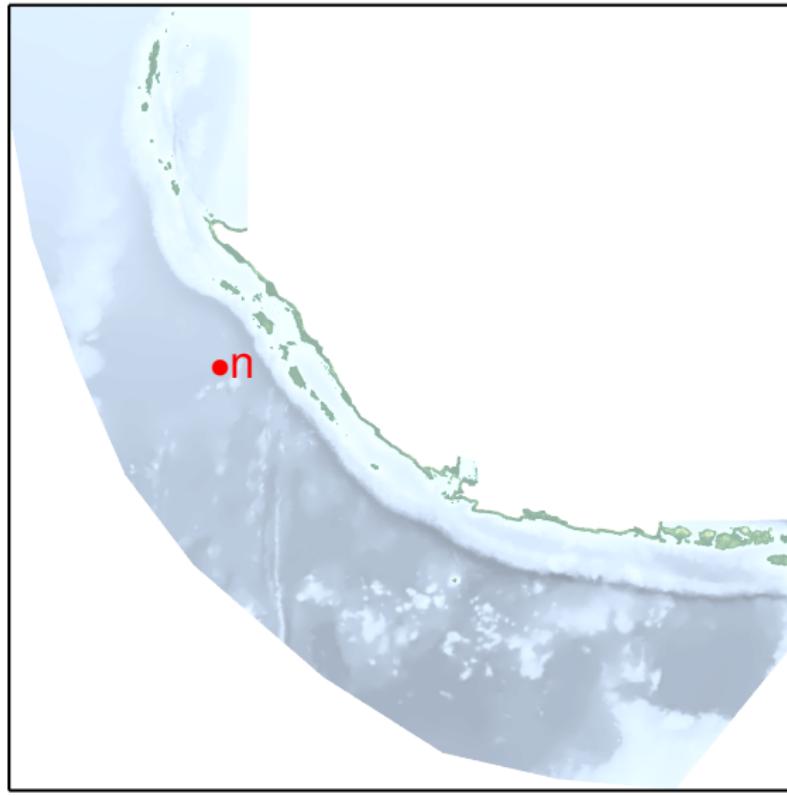
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But how to derive SFC for highly unstructured grid?



- Construct SFC like 3D Hilbert curve in particle code
Gadget-2 (communication with T. Rung, TU
Hamburg-Harburg)

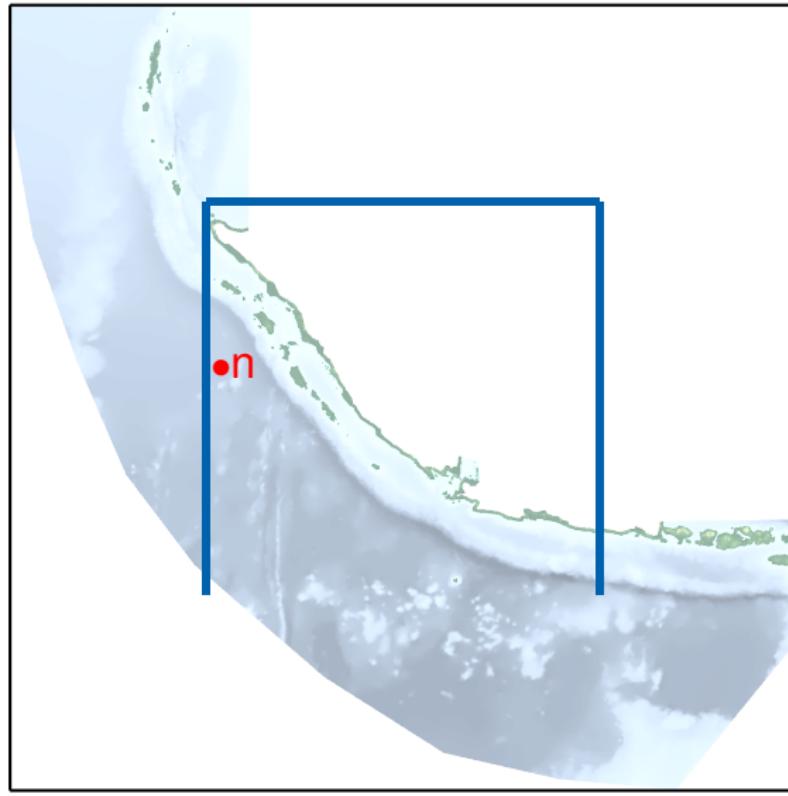
SFC construction



For all nodes n calculate the index in the Hilbert curve as a quad number:

$\text{SFC_index}(n) =$

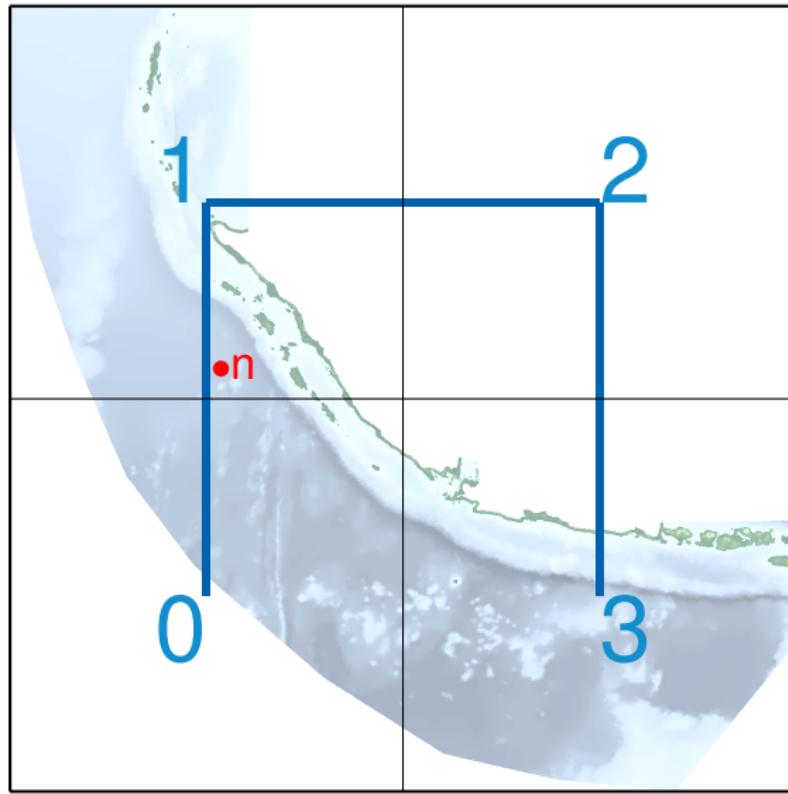
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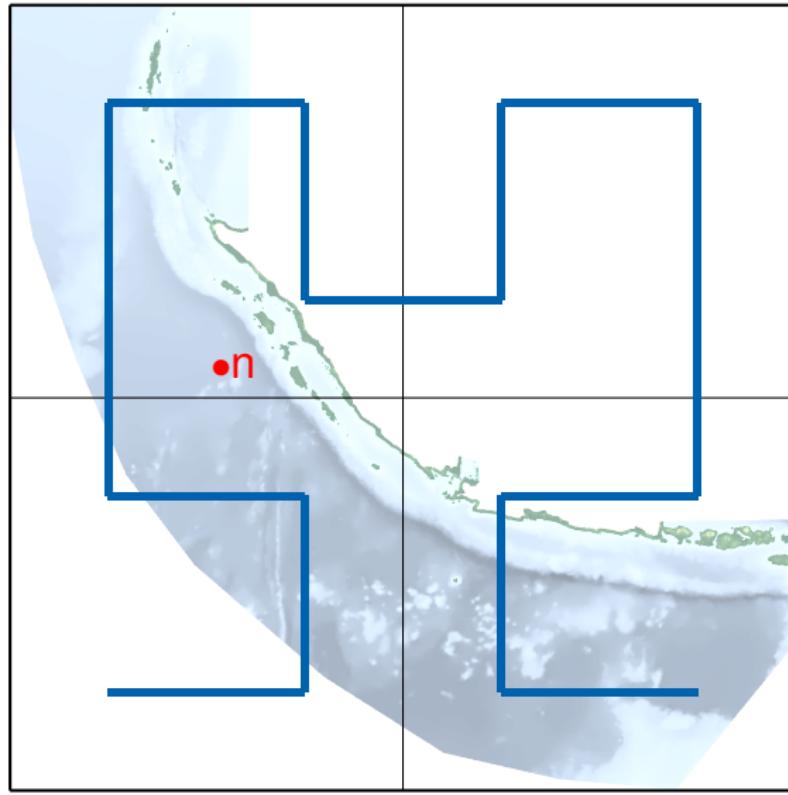


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1

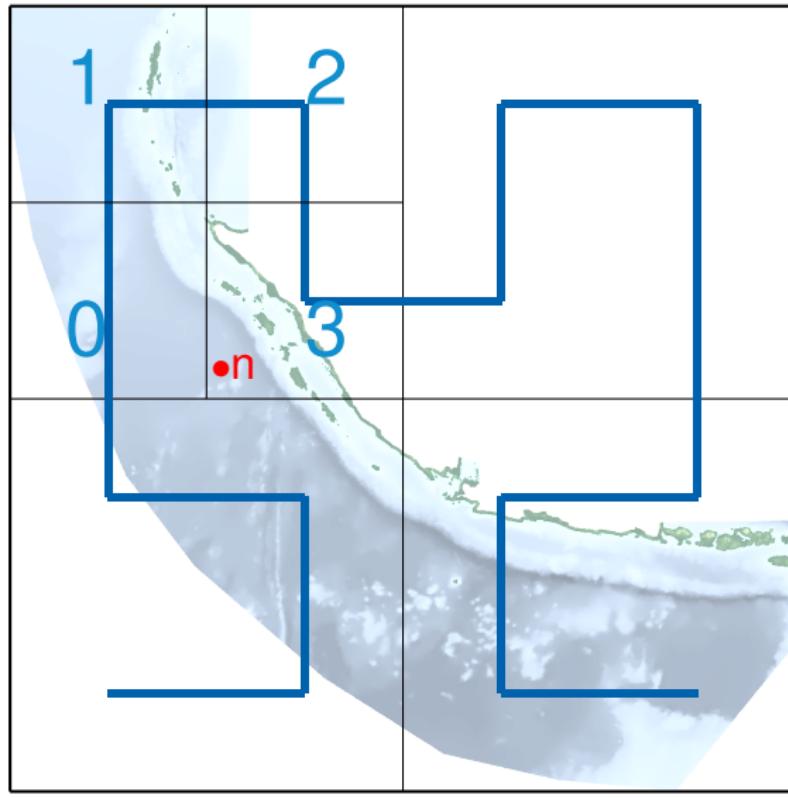
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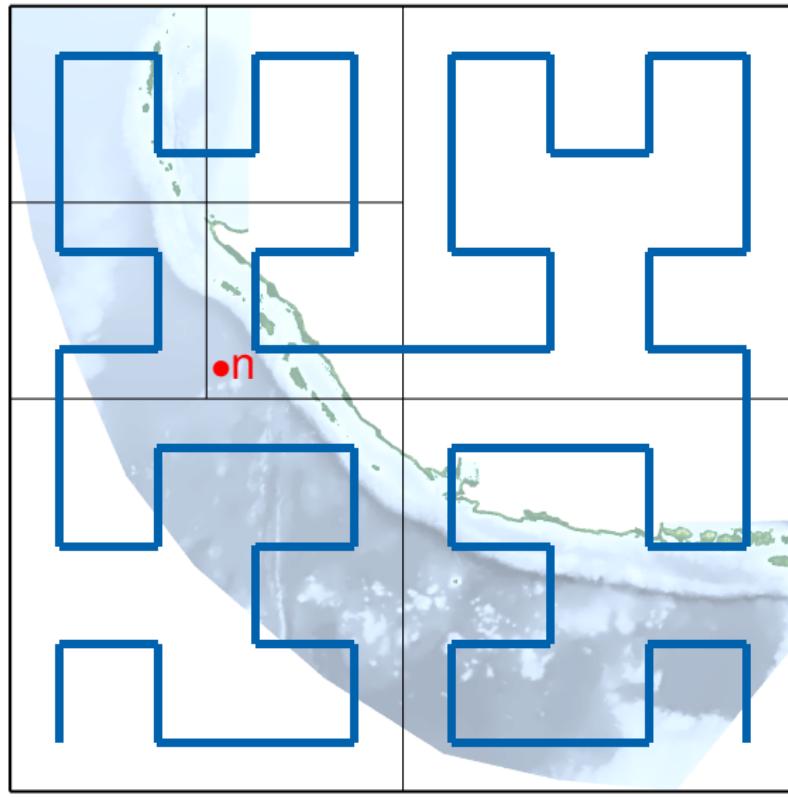
SFC construction



For all nodes n calculate the index in the Hilbert curve as a quad number:

$$\text{SFC_index}(n) = 13$$

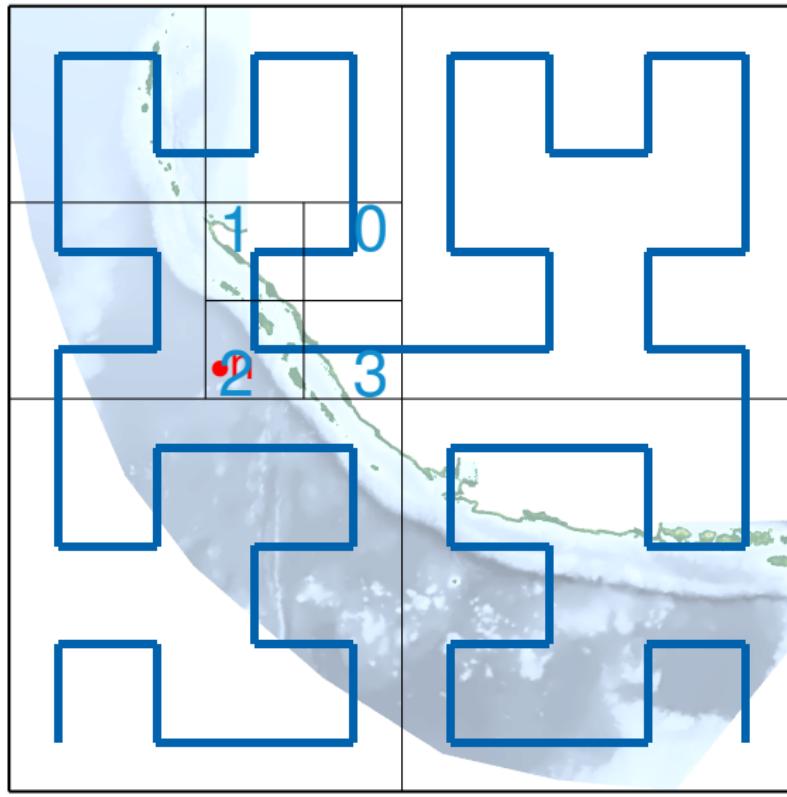
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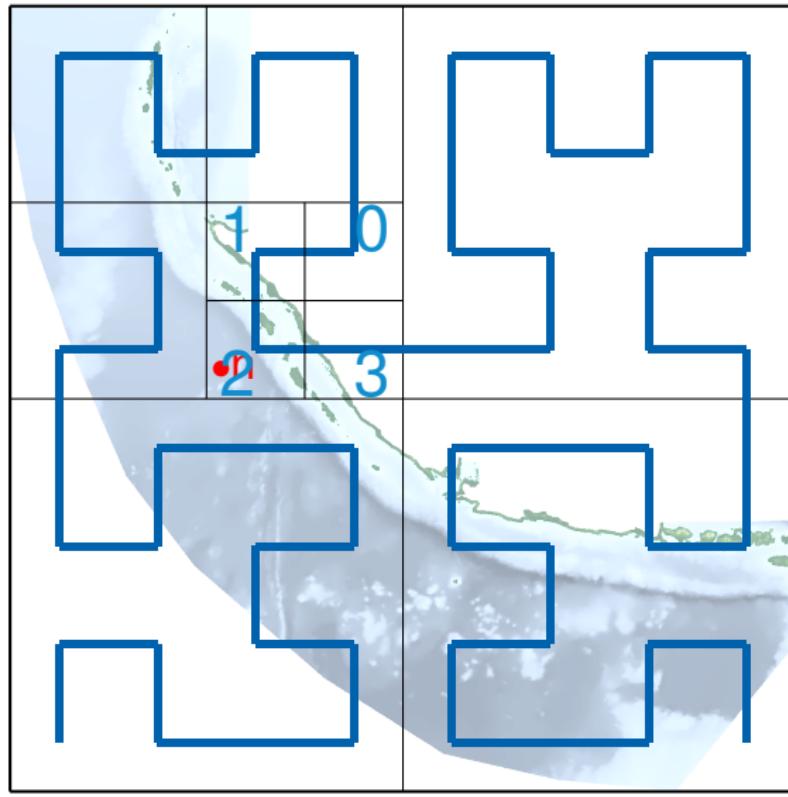
SFC construction



For all nodes n calculate the index in the Hilbert curve as a quad number:

$$\text{SFC_index}(n) = 132$$

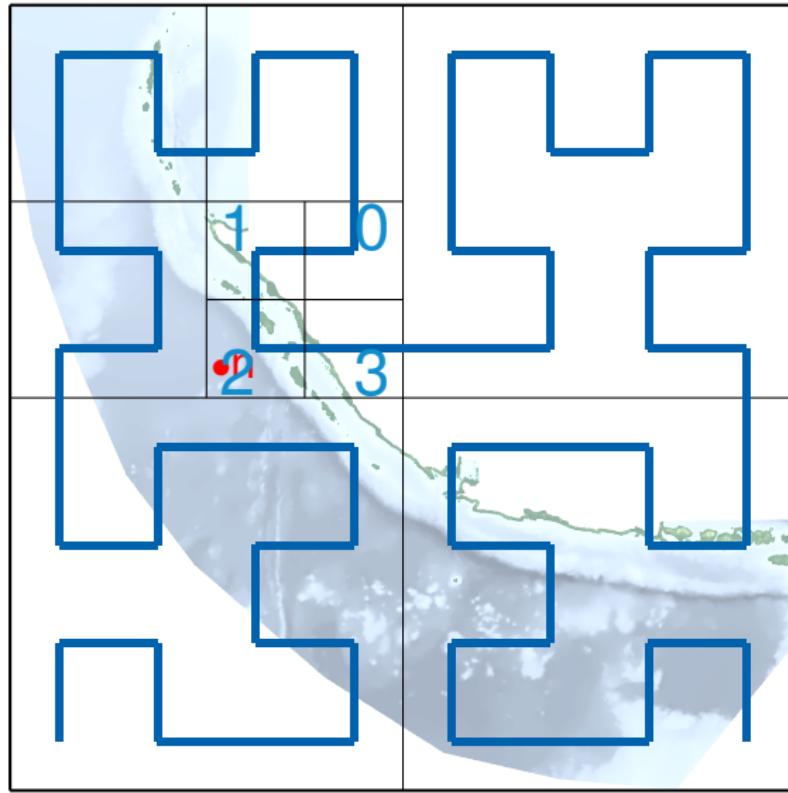
SFC construction



For all nodes n calculate the index in the Hilbert curve as a quad number:

$$\text{SFC_index}(n) = 132\dots$$

SFC construction



For all nodes n calculate the index in the Hilbert curve as a quad number:

$$\text{SFC_index}(n) = 132\dots$$

e.g. for 8 levels:

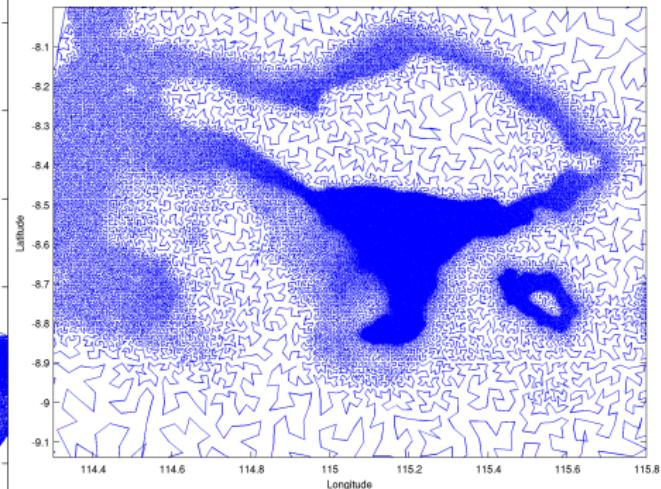
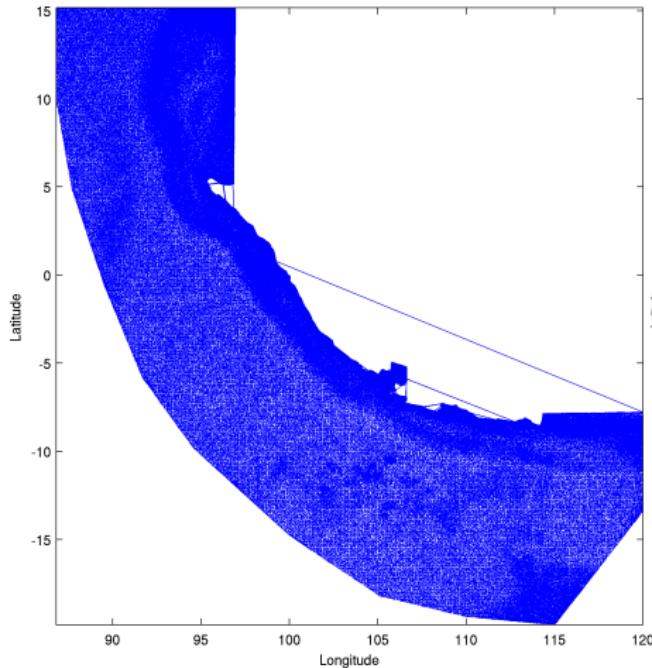
$$\text{SFC_index}(n) = 1 \cdot 4^8 + 3 \cdot 4^7 + 2 \cdot 4^6 + \dots$$

SFC reordering

- Reorder the nodes according to SFC_index.
- Reorder the elements
 - by an SFC separately, or
 - numerically by node indicees
(more efficient for TsunAWI)
- Edges are constructed in TsunAWI (sorted along the nodes)

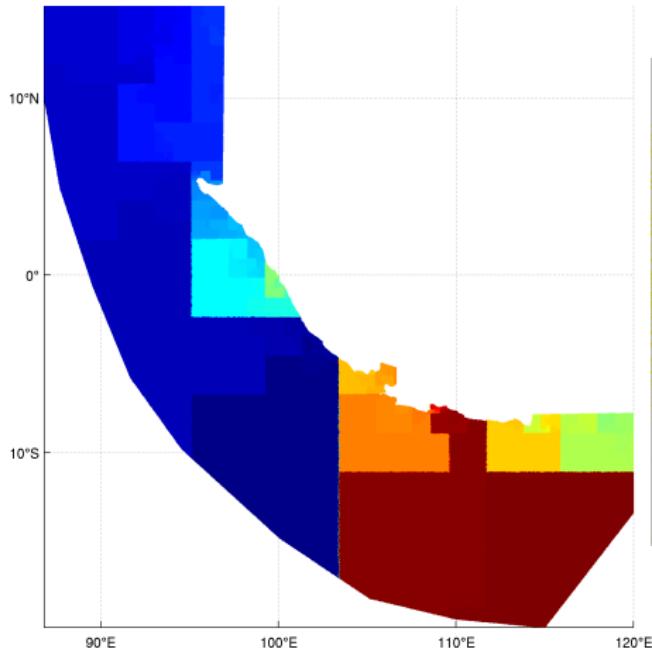
SFC ordering of the nodes

for TsunAWI regional indonesian grid

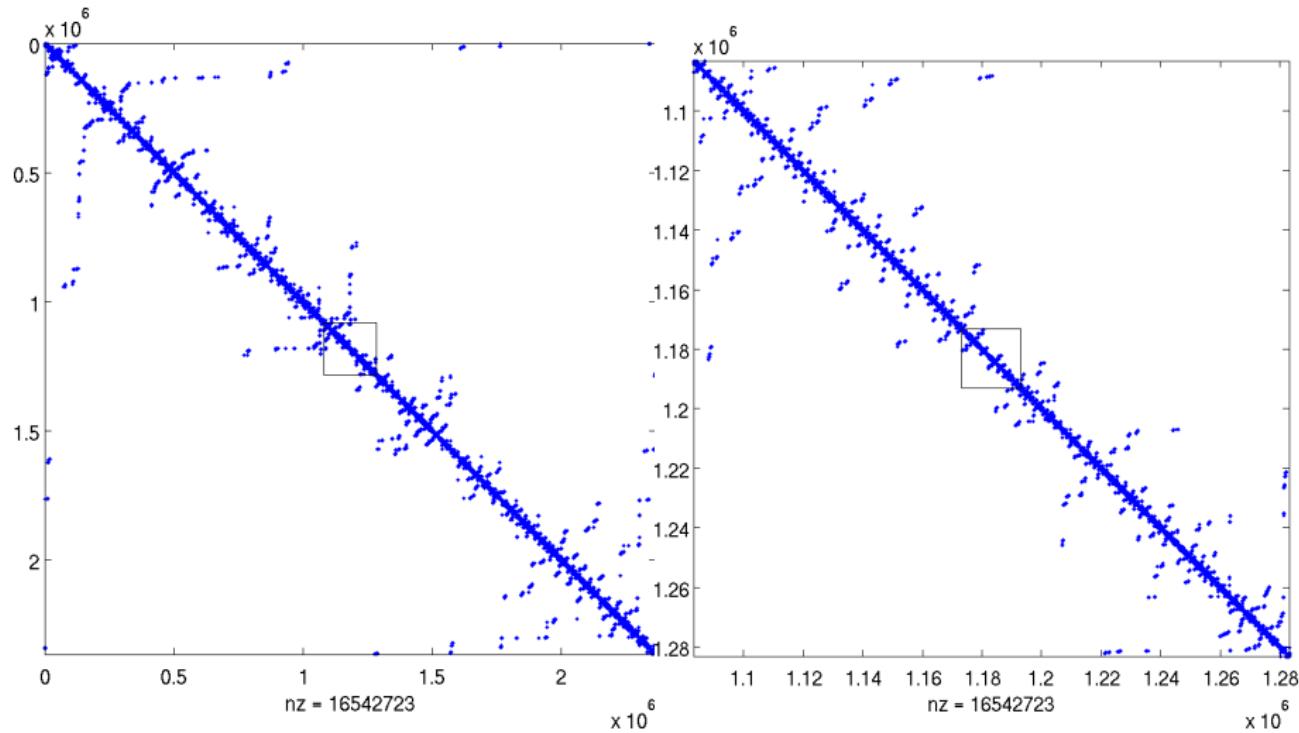


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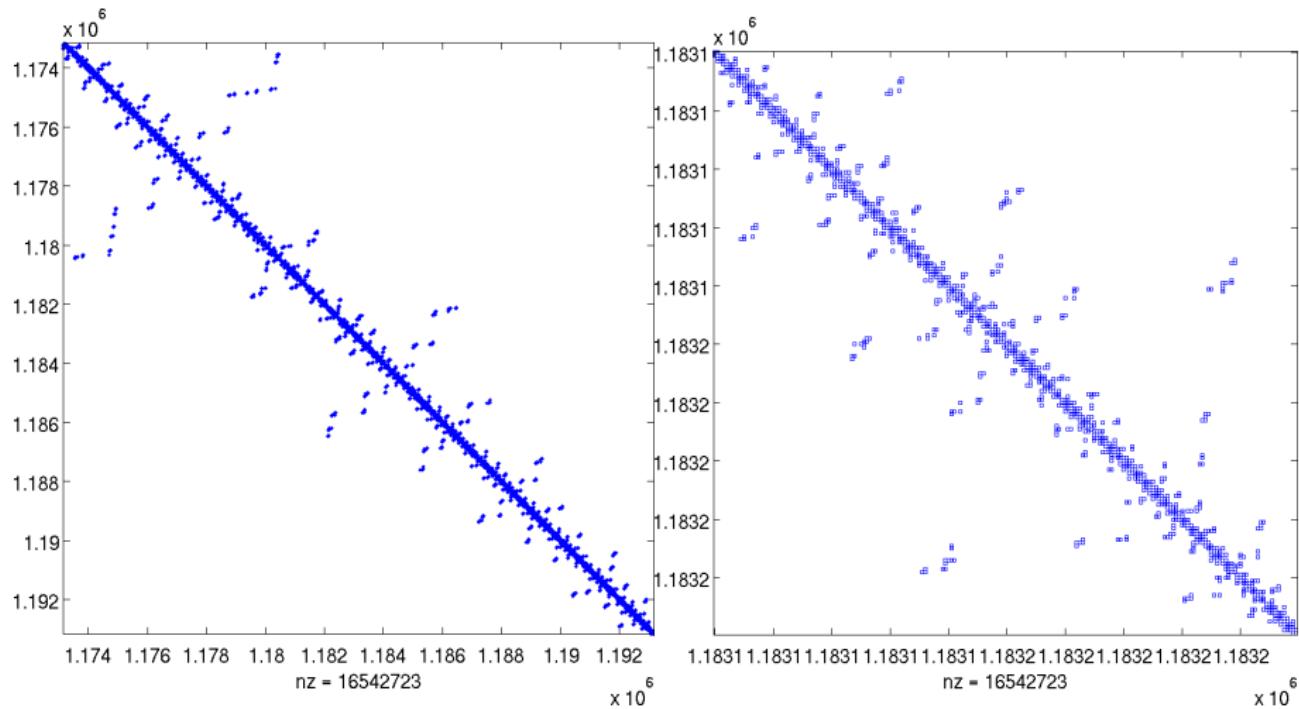
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adjacency matrix for SFC sorted grid

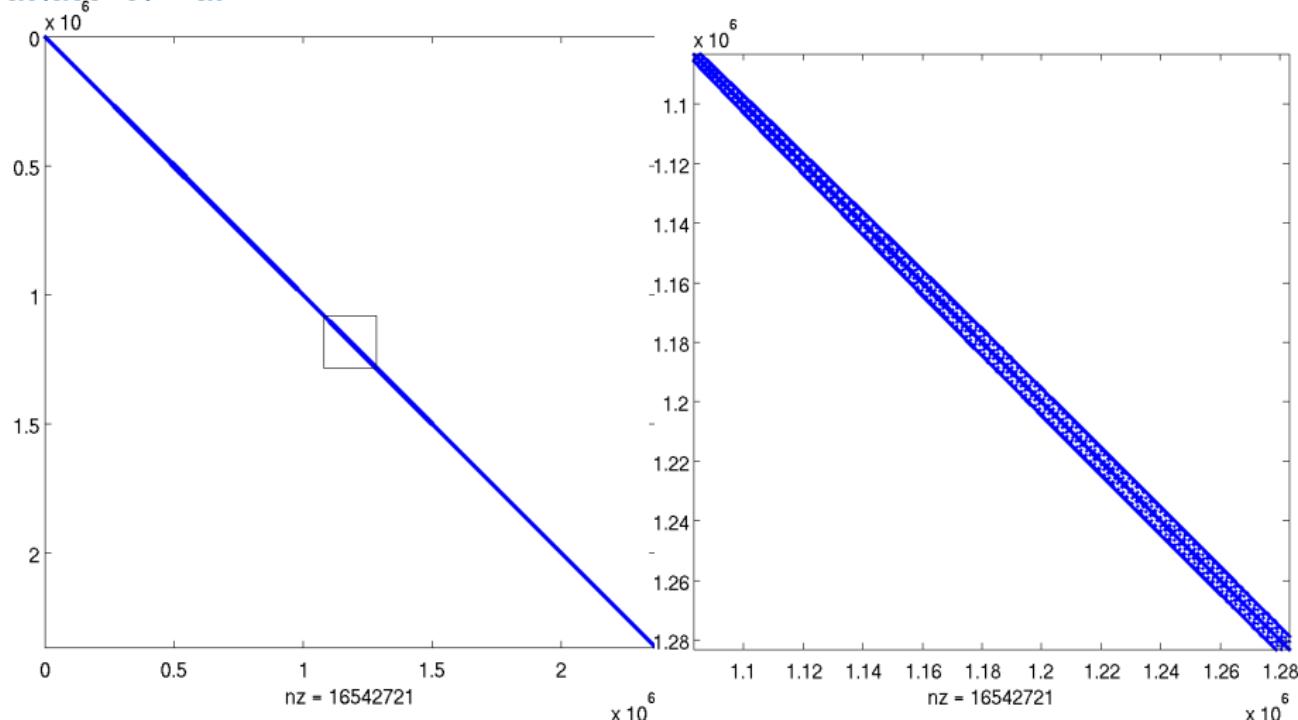


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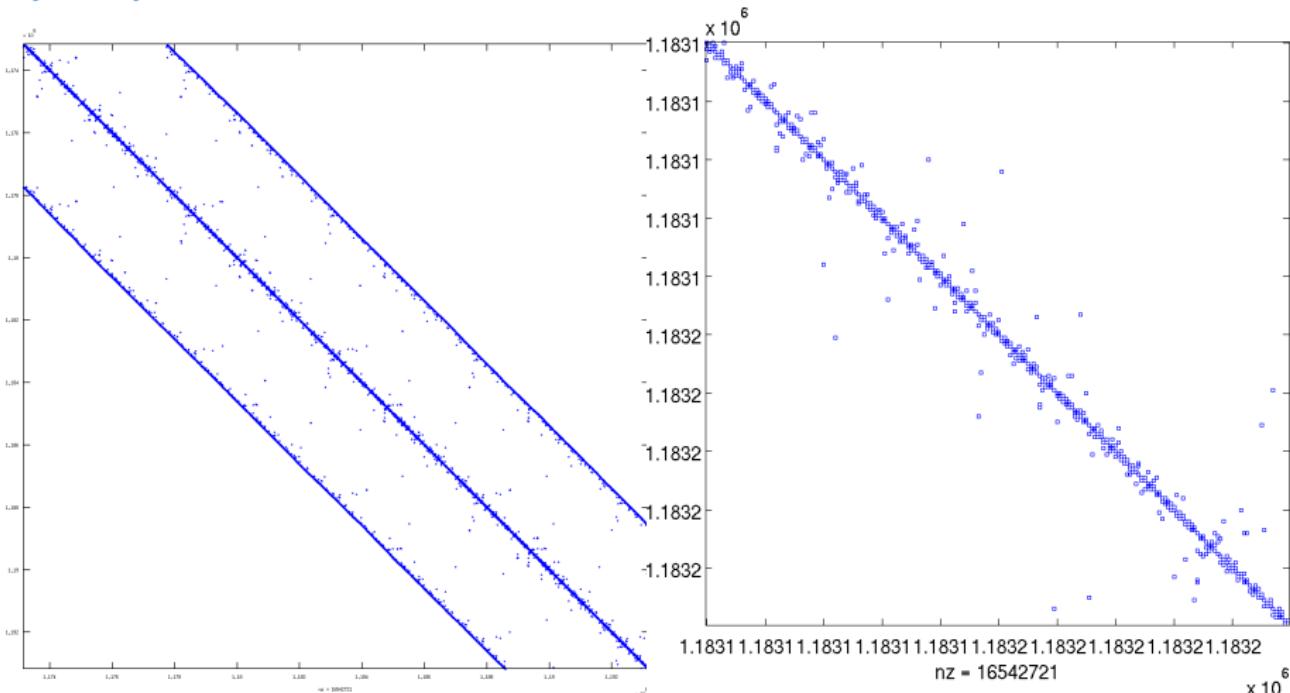
Comparison: RCM ordering

adjacency matrix



RCM (reverse Cuthill McKee) ordering obtained via adjacency matrix and Matlab symrcm for sparse matrices.

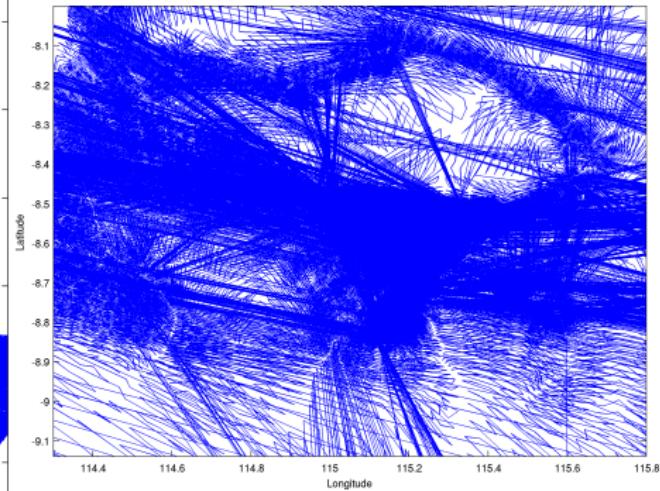
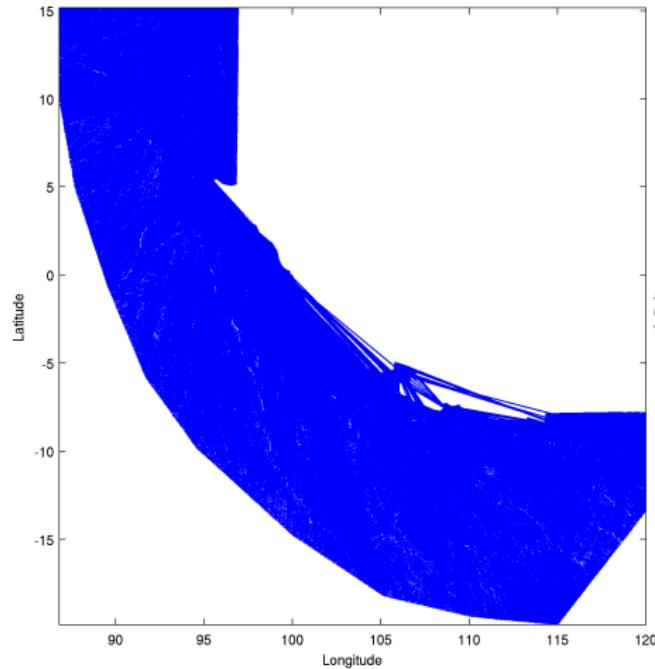
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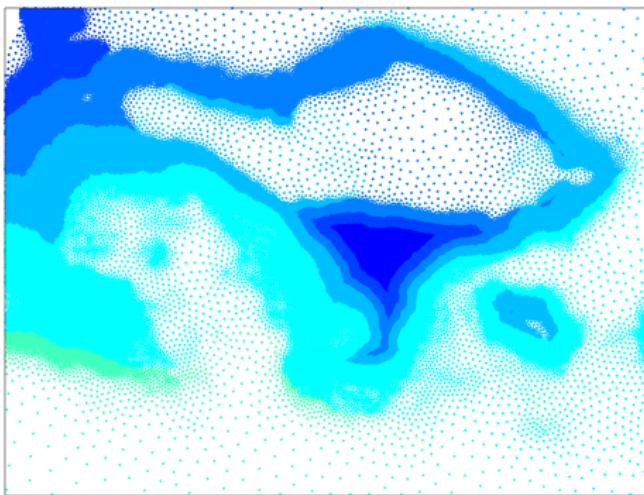
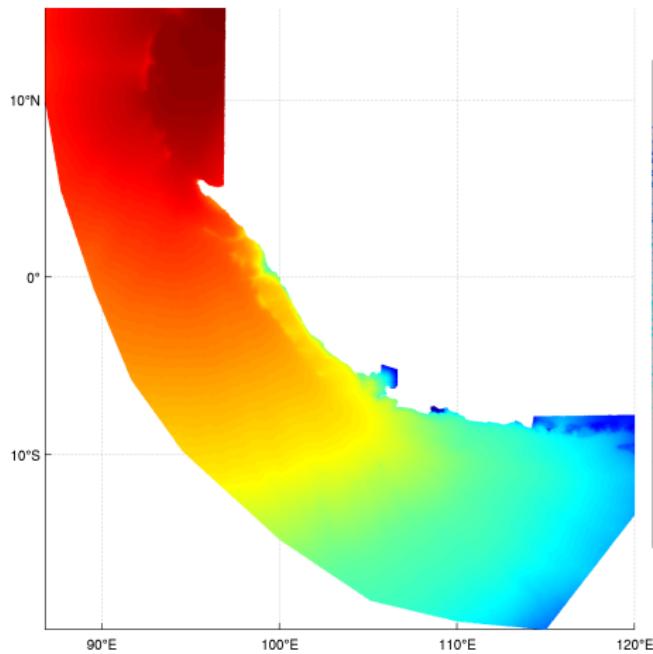
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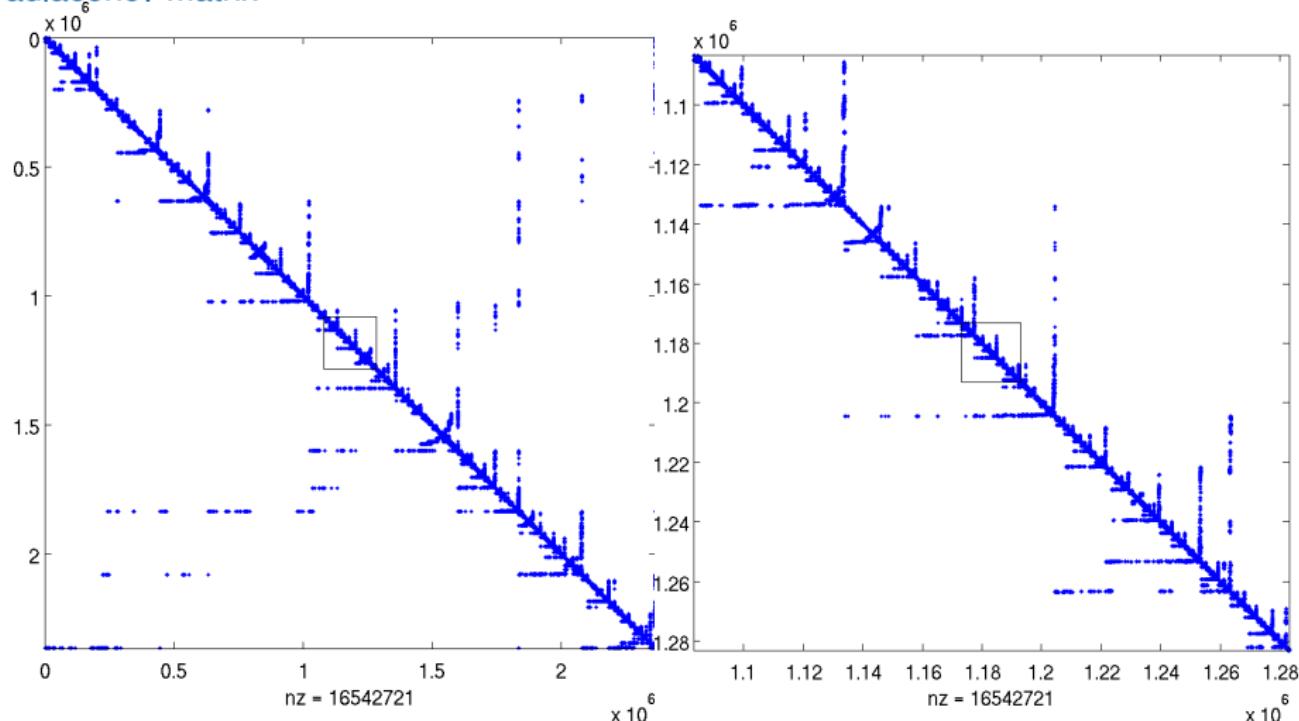
Comparison: RCM ordering

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Comparison: AMD ordering

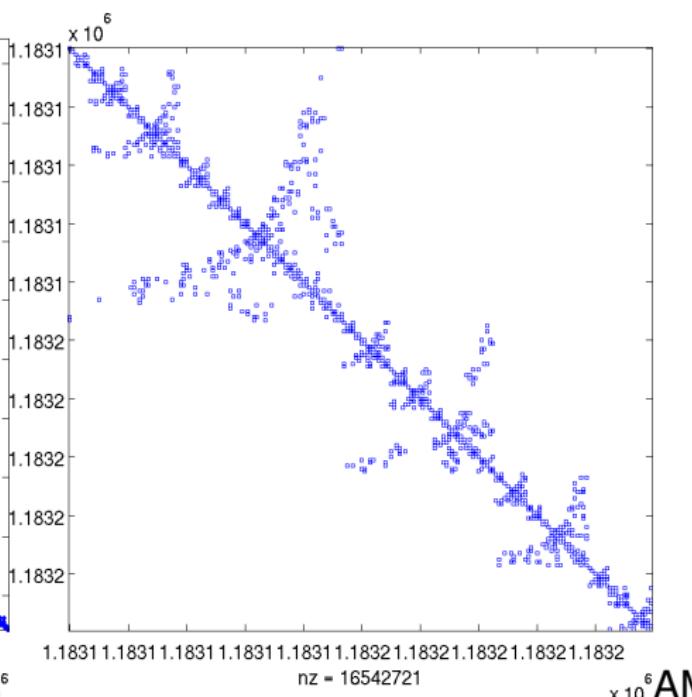
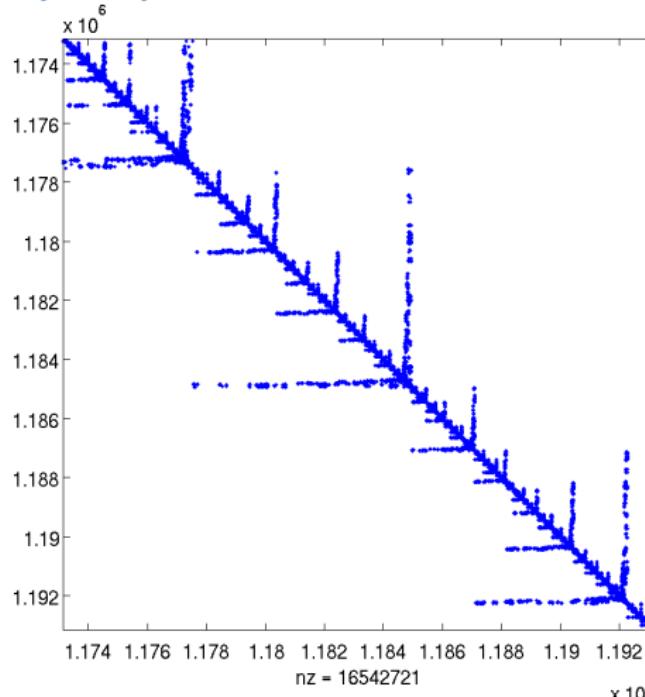
adjacency matrix



AMD (approximate minimum degree) ordering obtained via adjacency matrix and Matlab symamd for sparse matrices.

Comparison: AMD ordering

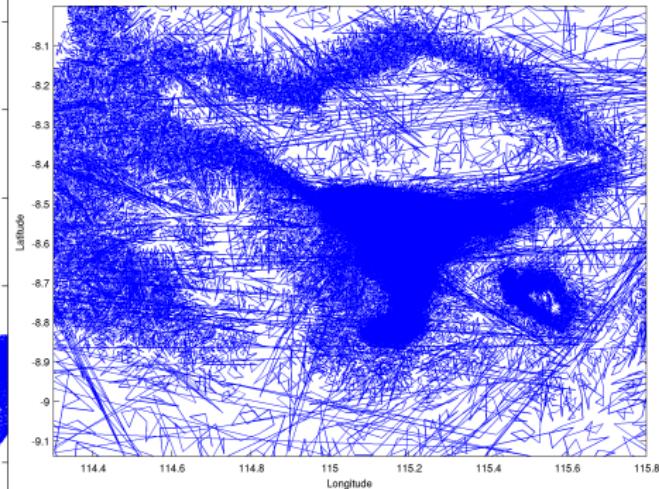
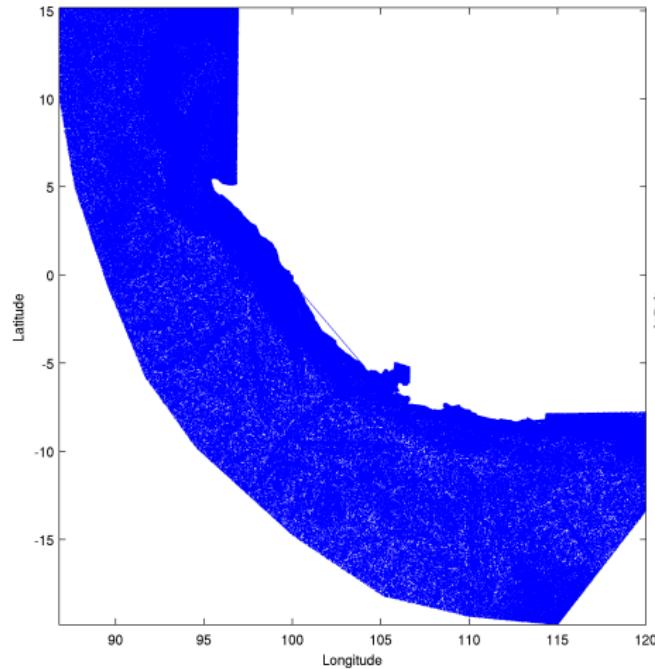
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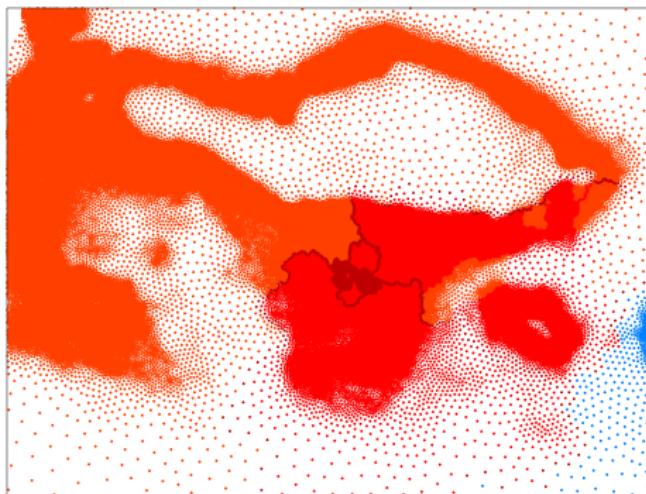
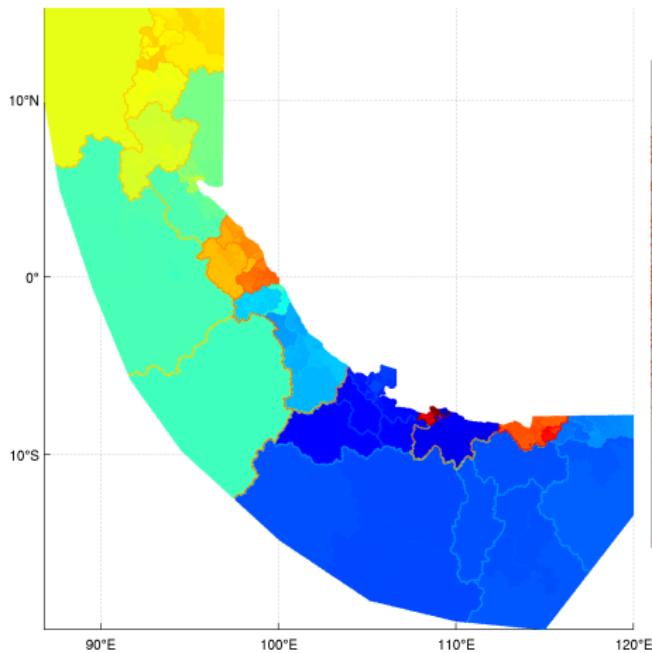
Comparison: AMD ordering

for TsunAWI regional indonesian grid



Comparison: AMD ordering

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SFC compared to unsorted, RCM, SymAMD

computation time: IBM Power6

Computational time [seconds] for timestep on a cluster node
1× IBM Power6 (4 Cores, 2× hyperthreading)

	OMP_NUM_THREADS			
	1	2	4	8
orig.	9.77	4.08	2.91	1.57
RCM	2.78	1.77	0.97	0.69
AMD	2.76	1.42	0.95	0.66
SFC	2.69	1.58	0.92	0.60

SFC compared to unsorted, RCM, SymAMD

Hardware counters: IBM Power6

IBM Hardware counter hpmcount for 1000 timesteps on
1× IBM Power6 (4 Cores, 2× hyperthreading,
OMP_NUM_THREADS=8)

	hpmcount event	
	L2 cache misses	Number of loads per load miss
orig.	274,478,564,540	17.8
RCM	57,244,100,260	64.0
AMD	54,709,662,295	65.6
SFC	49,980,798,689	88.5

SFC compared to unsorted, RCM, SymAMD

computation time: Intel Xeon Nehalem-EX

Computational time [seconds] for one timestep on
one blade SGI Altix UV (HLRN, ZIB Berlin and RRZN Hannover)
2× Intel Xeon 5570 (8 Cores, 2× hyperthreading)

	OMP_NUM_THREADS					
	1	2	4	8	16	32
orig.	3.84	2.16	1.48	0.89	0.52	0.40
RCM	1.64	1.12	0.59	0.35	0.20	0.19
AMD	1.47	0.77	0.50	0.30	0.18	0.16
SFC	1.47	0.90	0.51	0.31	0.17	0.14

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RCM	1.64	1.12	0.59	0.35	0.20	0.19	0.37	0.32
AMD	1.47	0.77	0.50	0.30	0.18	0.16	0.32	0.19
SFC	1.47	0.90	0.51	0.31	0.17	0.14	0.30	0.18

Remark on OpenMP

importance of first touch for data locality

```
allocate(array(dim))
```

```
array(:) = 0.
```

```
!$OMP PARALLEL DO
do n=1,dim
array(n) = ...
end do
!$OMP END PARALLEL DO
```

Remark on OpenMP

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```

properties of resorting by a SFC

SFC is a very valuable method, because

- it is cheap to compute
- provides good data locality
on **all** levels of the memory hierarchy
- as domain decomposition, it keeps interfaces small
(though not optimal)

work to do

- Influence of SFC ordering on
 - ILU based preconditioners
 - fill-in
 - computational load
 - convergence rate
 - sparse matrix computations in general
- SFC compared to generic partitioning algorithms
(MeTiS, scotch,...)
- TsunAWI
 - further optimize OpenMP parallelization
 - MPI parallelization